A NEW ALGORITHM FOR COMPUTING THE INVERSE AND GENERALIZED INVERSE OF THE SCALED FACTOR CIRCULANT MATRIX*

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Abstract

A new algorithm for finding the inverse of a nonsingular scaled factor circulant matrix is presented by the Euclid's algorithm. Extension is made to compute the group inverse and the Moore-Penrose inverse of the singular scaled factor circulant matrix. Numerical examples are presented to demonstrate the implementation of the proposed algorithm.

Mathematics subject classification: 15A21, 65F15. Key words: Scaled factor circulant matrix, Inverse, Group inverse, Moore-Penrose inverse.

1. Introduction

Circulant matrices, as an important class of special matrices, have a wide range of interesting applications [12–19]. They have in recent years been applied in many areas, see, e.g., [2, 3, 6, 10, 11, 15, 17]. Scaled circulant permutation matrices and the matrices that commute with them are natural extensions of this well-studied class, see, e.g., [1, 20–23]. In particular, it will be seen that *r*-circulant matrices [10, 11] are precisely those matrices commuting with the scaled circulant permutation matrix.

This paper presents an efficient algorithm to compute the inverse of a nonsingular scaled factor circulant matrix or to compute the group inverse and Moore-Penrose inverse of the circulant matrix when it is singular. The algorithm has small computational complexity. It is a notable character of the algorithm that the singularity of the scaled factor circulant matrix need not be priori known.

We define \mathcal{R} as the scaled circulant permutation matrix, that is,

$$\mathcal{R} = \begin{pmatrix} 0 & d_1 & 0 & \dots & 0 & 0 \\ 0 & 0 & d_2 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 0 & d_{n-1} \\ d_n & 0 & 0 & \dots & 0 & 0 \end{pmatrix}_{n \times n}$$
(1.1)

This paper deals with the case where \mathcal{R} is nonsingular ($d_i \neq 0$ and fixed).

It is easily verified that the polynomial $g(x) = x^n - d_1 d_2 \dots d_n$ is both the minimal polynomial and the characteristic polynomial of the matrix \mathcal{R} . In addition, \mathcal{R} is nondergatory.

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Moreover, \mathcal{R} is normal if and only if $|d_1| = |d_2| = \cdots = |d_n|$, where $|d_i|, i = 1, \cdots, n$ denote the modulus of the complex number $d_i, i = 1, \cdots, n$.

Definition 1.1. An $n \times n$ matrix A over \mathbb{C} is called a scaled factor circulant matrix if A commutes with \mathcal{R} , that is,

$$A\mathcal{R} = \mathcal{R}A,\tag{1.2}$$

where \mathcal{R} is given in (1.1).

Let $\mathcal{R}SFCM_n$ be the set of all complex $n \times n$ matrices which commute with \mathcal{R} . In the following, with $A = \operatorname{scacirc}_{\mathcal{R}}(a_0, a_1, \cdots, a_{n-1})$ we denote the scaled factor circulant matrix A whose first row is $(a_0, a_1, \ldots, a_{n-1})$. Remark that the first row of A completely defines the matrix. Indeed, since \mathcal{R} is nonderogatory, Eq. (1.2) is fulfilled if and only if $A = f(\mathcal{R})$ for some polynomial f. Furthermore, $\mathcal{R}SFCM_n$ is a vector space of dimension n, and there is a clear one-to-one correspondence between the polynomials of degree at most n-1 and the numbers a_0, \cdots, a_{n-1} .

For an $m \times n$ matrix A, any solution to the matrix equation AXA = A is called a *generalized* inverse of A. In addition, if X satisfies X = XAX, then A and X are said to be semi-inverses, see, e.g., [2].

In this paper we only consider square matrices A. In [8, p.51] the smallest positive integer k for which rank (A^{k+1}) =rank (A^k) holds is called the *index* of A. If A has index 1, the generalized inverse X of A is called the *group inverse* $A^{\#}$ of A. Clearly, A and X are group inverses if and only if they are semi-inverses and AX = XA.

In [4, 5] a semi-inverse X of A was considered in which the nonzero eigenvalues of X are the reciprocals of the nonzero eigenvalue of A. These matrices were called *spectral inverses*. It was shown in [5] that a nonzero matrix A has a unique spectral inverse, A^s , if and only if A has index 1: when A^s is the group inverse $A^{\#}$ of A.

2. The Properties of the Scaled Factor Circulant Matrix

Lemma 2.1. ([1]) If \mathcal{R} is a scaled circulant permutation matrix, and if k is a positive integer, then $\mathcal{R}^k = D^{(k)}C^k$, where $D^{(k)}$ is the diagonal matrix whose (j,j) entry is $\prod_{t=j}^{j+k-1} d_t$ for $1 \leq j \leq n$ and $C = circ(0, 1, 0, \dots, 0)$ is the circulant permutation. Furthermore,

$$\mathcal{R}^n = (\prod_{j=1}^n d_j) I_n, \quad \det \mathcal{R} = (-1)^{n-1} \prod_{j=1}^n d_j.$$

Let $\omega = \exp(\frac{2\pi i}{n})$ be a primitive *n*th root of unity. Then $\omega_j = d\omega^j$, $j = 0, 1, \dots, n-1$ are the distinct roots of g(x), where $g(x) = x^n - d_1 d_2 \cdots d_n$, and

$$d = (\prod_{t=1}^{n} d_t)^{\frac{1}{n}} \neq 0.$$
(2.1)

Let F be the $n \times n$ unitary Fourier matrix such that

$$F_{ij} = \frac{1}{\sqrt{n}} \omega^{(i-1)(j-1)} \quad \text{for } 1 \le i, \ j \le n.$$
(2.2)

Let

$$\Delta = \operatorname{diag}(\delta_1, \delta_2, \cdots, \delta_n), \tag{2.3}$$

where the elements δ_j of Δ are computed by the recursion formula

$$\delta_{j+1} = \frac{d}{d_j}\delta_j, \quad 1 \le j \le n, \quad \delta_{n+1} = \delta_1 = 1.$$

Lemma 2.2. ([1]) Let $A = \text{scacirc}_{\mathcal{R}}(a_0, a_1, \cdots, a_{n-1})$ be a scaled factor circulant matrix over the complex field \mathbb{C} . Then

$$\sigma(A) = \{\lambda_j | \lambda_j = f(d\omega^j) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} (d\omega^j)^i | 0 \le j \le n-1\}$$
(2.4)

is the spectrum of A and

$$A = f(\mathcal{R}) = a_0 I + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} \mathcal{R}^i,$$
(2.5)

where

$$f(x) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} x^i.$$
 (2.6)

The polynomial (2.6) will be called the representor of the scaled factor circulant matrix A.

Lemma 2.3. ([1]) Let $A = \operatorname{scacirc}_{\mathcal{R}}(a_0, a_1, \cdots, a_{n-1})$ be a scaled factor circulant matrix over the complex field \mathbb{C} . If F is the Fourier matrix, then

$$A = (\Delta F) \operatorname{diag}(\lambda_0, \cdots, \lambda_i, \cdots, \lambda_{n-1}) (\Delta F)^{-1}, \qquad (2.7)$$

where Δ is given by (2.3) and λ_j , $j = 0, 1, \dots, n-1$ are the eigenvalues of A given by (2.4).

Let D_n denote the multiplicative semigroup of all $n \times n$ diagonal complex matrices. By Lemma 1 in [2, p.27] the mapping

$$A \to (\triangle F)^{-1} A(\triangle F)$$

is a semigroup isomorphism of $\mathcal{R}SFCM_n$ onto D_n , where F and \triangle are defined by (2.2) and (2.3), respectively.

Let $A = \text{scacirc}_{\mathcal{R}}(a_0, a_1, \dots, a_{n-1}) \in \mathcal{R}SFCM_n$ be a scaled factor circulant matrix. Then $\sigma(A) = \{\lambda_i | i = 0, 1, \dots, n-1\}$ by (2.4). Let

$$T_i = \begin{cases} 0, & \text{if } \lambda_i = 0, \\ 1/\lambda_i, & \text{if } \lambda_i \neq 0, \end{cases}$$

for $i = 0, 1, \dots, n - 1$. If

$$B = (\triangle F) \operatorname{diag}(T_0, \cdots, T_i, \cdots, T_{n-1}) (\triangle F)^{-1},$$

then by Theorem 1 of [2], $B = A^s$, the spectral inverse of A.

Since each A in $\mathcal{R}SFCM_n$ has index 1, A^s is also the group inverse $A^{\#}$ of A. Moreover, if \mathcal{R} is normal, then by Theorem 1 of [2], $A^s = A^+$, where A^+ denotes the Moore-Penrose inverse of A.

We summarize the above discussions in the following theorems.

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Theorem 2.1. Let $A \in M_n$. Then $A \in \mathcal{R}SFCM_n$ if and only if $(\triangle F)^{-1}A(\triangle F)$ is a diagonal matrix. Let $A \in \mathcal{R}SFCM_n$. If A is a singular matrix, then $A^s = A^{\#} \in \mathcal{R}SFCM_n$. If \mathcal{R} is normal, then $A^+ \in \mathcal{R}SFCM_n$ and $A^+ = A^{\#}$.

Theorem 2.2. If $A = \text{scacirc}_{\mathcal{R}}(a_0, a_1, \dots, a_{n-1})$ is nonsingular, then $f(\omega_i) \neq 0$, where

$$f(x) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} x^i, \quad \omega_j = d\omega^j, \quad j = 0, 1, \cdots, n-1$$

are the distinct roots of g(x). If A is singular and has k zero eigenvalues, then there are $\omega_{i_0}, \omega_{i_1}, \dots, \omega_{i_{k-1}}$, such that $f(\omega_{i_j}) = 0$, for $j = 0, 1, \dots, k-1$. Conversely, if there exists ω_k satisfying $f(\omega_k) = 0$, then the scaled factor circulant matrix A is singular.

Proof. According to Theorem 2.1, we know that $(\triangle F)^{-1}A(\triangle F) = D$, where

$$D = \operatorname{diag}(f(\omega_0), f(\omega_1), \cdots, f(\omega_{n-1})),$$

and $\omega_j = d\omega^j$, $j = 0, 1, \dots, n-1$ are the distinct roots of g(x). Thus $A \triangle F = \triangle FD$. Since $\triangle F$ is a nonsingular matrix, then

$$\operatorname{rank} A = \operatorname{rank} A \triangle F = \operatorname{rank} \triangle F D = \operatorname{rank} D.$$

If there exist $\omega_{i_0}, \omega_{i_1}, \cdots, \omega_{i_{k-1}}$ such that $f(\omega_{i_j}) = 0$, for $j = 0, 1, \cdots, k-1$, then there are $\omega_{i_k}, \omega_{i_{k+1}}, \cdots, \omega_{i_{n-1}}$ such that $f(\omega_{i_j}) \neq 0$, for $j = k, k+1, \cdots, n-1$. Thus rank A = n-k.

Conversely, if rank A = n - k, then there exist $\omega_{i_k}, \omega_{i_{k+1}}, \cdots, \omega_{i_{n-1}}$ such that $f(\omega_{i_j}) \neq 0$, for $j = k, k + 1, \cdots, n - 1$. Therefore, there are $\omega_{i_0}, \omega_{i_1}, \cdots, \omega_{i_{k-1}}$ such that $f(\omega_{i_j}) = 0$, for $j = 0, 1, \cdots, k - 1$.

In addition, let $A, B \in \mathcal{R}SFCM_n$. Then $AB = BA \in \mathcal{R}SFCM_n$. If A is a nonsingular matrix, then $A^{-1} \in \mathcal{R}SFCM_n$. Thus $\mathcal{R}SFCM_n$ is a ring.

Polynomial ring has an intimate relation to the scaled factor circulant matrix ring. Let P(x) be the polynomial ring. For all f(x) in P(x), the degree of f(x) is denoted by $\deg(f(x))$. Let $P_{n-1}(x)$ be the quotient ring $P(x)/\langle x^n - d_1d_2\cdots d_n\rangle$, where $\langle x^n - d_1d_2\cdots d_n\rangle$ is an ideal. Define φ as a function which maps scaled factor circulant matrix ring onto the polynomial ring by

$$\varphi(A) \mapsto f(x) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} x^i,$$

where $A = \operatorname{scacirc}_{\mathcal{R}}(a_0, a_1, \cdots, a_{n-1}).$

Then, we can conclude that φ is a ring isomorphism. The scaled factor circulant matrix ring and the polynomial quotient ring $P_{n-1}(x)$ are isomorphic. So, if A is nonsingular, then φ maps the inverse of A onto the inverse of the representor f(x) of A.

3. Main Results

Theorem 3.1. Let $A = \text{scacirc}_{\mathcal{R}}(a_0, a_1, \dots, a_{n-1})$ be a scaled factor circulant matrix which is nonsingular, with the representor of A being

$$f(x) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} x^i.$$

Then there exists a polynomial

$$u(x) = b_0 + \sum_{i=1}^{n-1} b_i (\prod_{t=1}^i d_t)^{-1} x^i$$

such that $u(\omega_j) = 1/f(\omega_j)$, where $\omega_j, j = 0, 1, \dots, n-1$, are the roots of $g(x) = x^n - d_1 d_2 \cdots d_n$ and the inverse of A is given by

$$B = \operatorname{scacirc}_{\mathcal{R}}(b_0, b_1, \cdots, b_{n-1}).$$

Proof. From Theorem 2.2, we know that $f(x) = f(\omega_j) \neq 0, j = 0, 1, \dots, n-1$. Let

$$g(x) = \prod_{j=0}^{n-1} (x - \omega_j) = x^n - d_1 d_2 \cdots d_n.$$

Then f(x) and g(x) are coprime. Hence there exist u'(x) and v(x) satisfying

$$f(x)u'(x) + g(x)v(x) = 1.$$

When $x = \omega_j$, $j = 0, 1, \dots, n-1$, then g(x) = 0. Consequently, $f(\omega_j)u'(\omega_j) = 1$. Let

$$u(x) = u'(x) \operatorname{mod}(x^n - d_1 d_2 \cdots d_n).$$

Then deg(u(x)) < n. Since $\omega_j^n - d_1 d_2 \cdots d_n = 0$, and $u(\omega_j) = u'(\omega_j), j = 0, 1, \cdots, n-1$, the existence of u(x) in Theorem 3.1 is then proved.

For the scaled factor circulant matrix B we have

$$B = \operatorname{scacirc}_{\mathcal{R}}(b_0, b_1, \cdots, b_{n-1})$$

= $\triangle F \operatorname{diag}(u(\omega_0), u(\omega_1), \cdots, u(\omega_{n-1}))(\triangle F)^{-1}$
= $\triangle F \operatorname{diag}(1/f(\omega_0), 1/f(\omega_1), \cdots, 1/f(\omega_{n-1}))(\triangle F)^{-1}.$

Consequently, BA = I. Therefore, u(x) is the inverse of f(x) in the quotient ring $P_{n-1}(x)$. The polynomial u'(x) can be obtained by Euclid's Algorithm. This is the main idea of the algorithm for computing the inverse of the scaled factor circulant matrix.

To reduce the computation, suppose a is the leading coefficient of f(x) and $a \neq 0$, let f'(x) = f(x)/a. Then f(x) = af'(x). The leading coefficient of f'(x) is 1.

Theorem 3.2. Let $A = \text{scacirc}_{\mathcal{R}}(a_0, a_1, \dots, a_{n-1})$ be a singular scaled factor circulant matrix with the representor

$$f(x) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} x^i.$$

Suppose A has m nonzero eigenvalues. Without loss of generality, suppose $f(\omega_j) = 0$, for $j = m, m + 1, \dots, n - 1$, where $\omega_j, j = 0, 1, \dots, n - 1$, are roots of $g(x) = x^n - d_1 d_2 \cdots d_n$.

Let

$$g_1(x) = \prod_{j=0}^{m-1} (x - \omega_j), \quad g_2(x) = \prod_{j=m}^{n-1} (x - \omega_j), \quad f_1(x) = f(x)g_2(x).$$

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Then there exists a polynomial

$$u_1(x) = b'_0 + \sum_{i=1}^{n-1} b'_i (\prod_{t=1}^i d_t)^{-1} x^i$$

such that $u_1(\omega_j) = 1/f_1(\omega_j), j = 0, 1, \dots, m-1$. Let

$$u(x) = u_1(x)g_2(x) = b_0 + \sum_{i=1}^{n-1} b_i(\prod_{t=1}^i d_t)^{-1}x^i.$$

Then $B = \operatorname{scacirc}_{\mathcal{R}}(b_0, b_1, \cdots, b_{n-1})$ is the group inverse $A^{\#}$ of A. If \mathcal{R} is normal, then $B = \operatorname{scacirc}_{\mathcal{R}}(b_0, b_1, \cdots, b_{n-1})$ is the Moore-Penrose inverse A^+ of A.

Proof. Since

$$x^{n} - d_{1}d_{2}\cdots d_{n} = \prod_{j=0}^{n-1} (x - \omega_{j}),$$

it follows that $g_1(x)$ and $g_2(x)$ are coprime. From the condition of Theorem 3.2, we know that $g_1(x)$ and f(x) are coprime. So $f_1(x)$ and $g_1(x)$ are coprime, and there exist $u_2(x)$ and v(x) satisfying

$$f_1(x)u_2(x) + g_1(x)v(x) = 1.$$

When $x = \omega_j$, $j = 0, 1, \dots, m-1$, $g_1(x) = 0$, thus $f_1(\omega_j)u_2(\omega_j) = 1$. Let $u_1(x) = u_2(x) \mod(g_1(x))$. Then the existence of the $u_1(x)$ in Theorem 3.2 has been proved.

Since $u(x) = u_1(x)g_2(x)$, when $j = m, m+1, \dots, n-1, u(\omega_j) = 0$, when $j = 0, 1, \dots, m-1$,

$$u(\omega_j) = u_1(\omega_j)g_2(\omega_j) = g_2(\omega_j)/f_1(\omega_j) = 1/f(\omega_j).$$

The scaled factor circulant matrix B is given by

$$B = \operatorname{scacirc}_{\mathcal{R}}(b_0, b_1, \cdots, b_{n-1})$$

= $\triangle F \operatorname{diag}(u(\omega_0), u(\omega_1), \cdots, u(\omega_{n-1}))(\triangle F)^{-1}$
= $\triangle F \operatorname{diag}(1/f(\omega_0), 1/f(\omega_1), \cdots, 1/f(\omega_{m-1}), 0, \cdots, 0)(\triangle F)^{-1}.$

It follows from Theorem 2.1 that B is the group inverse $A^{\#}$ of A. If \mathcal{R} is normal, then B is the Moore-Penrose inverse A^+ of A.

Theorem 3.2 implies that for computing the group inverse $A^{\#}$ and the Moore-Penrose inverse A^+ of the singular scaled factor circulant matrix A, we only need to invert $f(x)g_2(x)$ in the quotient ring $P_{n-1}(x)/\langle g_1(x)\rangle$.

It can be verified that $g_2(x)$ is the largest common factor of f(x) and $g(x) = x^n - d_1 d_2 \cdots d_n$. In our computations, if $\deg(f_1(x)) > \deg(g_1(x))$, we can do polynomial division $f_1(x) = g_1(x)s(x) + f_{12}(x)$. As $f_{12}(\omega_j) = f_1(\omega_j)$, $j = 0, 1, \cdots, m-1$, $f_1(x)$ can be taken the place by $f_{12}(x)$.

A similar device was used in [24] for computing the inverses and the group inverses of FLS r-circulant matrices.

4. Inverting the Scaled Factor Circulant Matrix

The problem becomes how to evaluate u(x), v(x) when f(x), g(x) are known and satisfy f(x)u(x) + g(x)v(x) = 1. Using Euclid's algorithm:

$$g(x) = q_0(x)f(x) + r_1(x),$$

$$f(x) = q_1(x)r_1(x) + r_2(x),$$

$$r_1(x) = q_2(x)r_2(x) + r_3(x),$$

.....

$$r_{i-1}(x) = q_i(x)r_i(x) + r_{i+1}(x),$$

.....

Let $v_1(x) = 1$, $u_1(x) = -q_0(x)$, then $r_1(x) = f(x)u_1(x) + g(x)v_1(x)$. It is obvious that

$$r_2(x) = f(x) - q_1(x)[g(x) - q_0(x)f(x)]$$

= $[1 + q_0(x)q_1(x)]f(x) - g(x)q_1(x)$

Let $v_2(x) = -q_1(x), u_2(x) = 1 + q_0(x)q_1(x)$. We then have

$$r_2(x) = f(x)u_2(x) + g(x)v_2(x)$$

Suppose

$$r_j(x) = f(x)u_j(x) + g(x)v_j(x),$$

and that $v_j(x), u_j(x)$ have been computed when $j = 0, 1, \dots, i$. Then

$$\begin{aligned} r_{i+1}(x) &= r_{i-1}(x) - q_i(x)r_i(x) \\ &= f(x)u_{i-1}(x) + g(x)v_{i-1}(x) - q_i(x)[f(x)u_i(x) + g(x)v_i(x)] \\ &= f(x)[u_{i-1}(x) - q_i(x)u_i(x)] + g(x)[v_{i-1}(x) - q_i(x)v_i(x)]. \end{aligned}$$

Let $r_{i+1}(x) = f(x)u_{i+1}(x) + g(x)v_{i+1}(x)$. Then

$$u_{i+1}(x) = u_{i-1}(x) - q_i(x)u_i(x)$$

Now, we have obtained the recurrence formula for $u_i(x)$. So, computing u(x) is equivalent to computing a sequence of polynomials

$$q_0(x), r_1(x), u_1(x), \cdots, q_i(x), r_{i+1}(x), u_{i+1}(x), \cdots$$

If $r_{j+1} = \text{const}$, then $u(x) = \frac{u_{j+1}(x)}{r_{j+1}(x)}$. We can improve the method. If the division of polynomial becomes

$$r_{i-1}(x) = q_i(x)r_i(x) + c_{i+1}r_{i+1}(x),$$

where c_{i+1} is a nonzero number, then the recurrence formula becomes

$$u_{i+1}(x) = \frac{u_{i-1}(x) - q_i(x)u_i(x)}{c_{i+1}}.$$

Thus, we can make the leading coefficient of $r_i(x)$ equal to 1 by suitably choosing c_i .

Algorithm 4.1. Given a scaled factor circulant matrix

$$A = \operatorname{scacirc}_{\mathcal{R}}(a_0, a_1, \cdots, a_{n-1}),$$

the algorithm computes a scaled factor circulant matrix

$$B = \operatorname{scacirc}_{\mathcal{R}}(b_0, b_1, \cdots, b_{n-1}).$$

When A is nonsingular, B is the inverse of A. When A is singular, B is the group inverse of A. If \mathcal{R} is normal, then B is the Moore-Penrose inverse A^+ of the scaled factor circulant matrix A. Let

$$f(x) = a_0 + \sum_{i=1}^{n-1} a_i (\prod_{t=1}^i d_t)^{-1} x^i, \quad g(x) = x^n - d_1 d_2 \cdots d_n,$$

$$r_{-1}(x) = g(x), \quad r_0(x) = f(x), \quad u_{-1}(x) = 0, \quad u_0(x) = 1.$$

Perform the polynomial division with remainder

do
$$\begin{cases} r_{i-1}(x) = q_i(x)r_i(x) + r_{i+1}(x), \\ (\text{let } c_{i+1} \text{ be the leading coefficient of } r_{i+1}(x)), \\ r_{i+1}(x) \leftarrow r_{i+1}(x)/c_{i+1}, \\ u_{i+1}(x) \leftarrow [u_{i-1}(x) - q_i(x)u_i(x)]/c_{i+1}, i = 0, 1, \cdots \end{cases}$$
(4.1)

until $r_m(x) = 1$ or $r_m(x) = 0$. If $r_m(x) = 1$, then $u_m(x)$ is the representor of B. Then $B = u_m(\mathcal{R})$ is the inverse of A. If $r_m(x) = 0$, then $r_{m-1}(x)$ is the largest common factor of f(x) and g(x). Let $r(x) = r_{m-1}(x), r_{-1}(x) = g(x)/r_{m-1}(x), r_0(x) = f(x)r_{m-1}(x) \mod(r_{-1}(x)), u_{-1}(x) = 0, u_0(x) = 1$, go to (4.1). Now, if $r_{m'}(x) = 1$, then $u(x) = u_{m'}(x)r(x) \mod(g(x))$ is the representor of B. Thus $B = u(\mathcal{R})$ is the group inverse $A^{\#}$ of A. Moreover, if \mathcal{R} is normal, then $B = u(\mathcal{R})$ is the Moore-Penrose inverse A^+ of A.

5. Computational Complexity

If the matrix is nonsingular, the computational complexity is divided into two parts. Suppose that the order of the scaled factor circulant matrix A is n, deg (f(x)) = n - 1. First we discuss the computational complexity on the division of polynomials.

$$r_{i-1}(x) = q_i(x)r_i(x) + r_{i+1}(x), r_{i+1}(x) \leftarrow r_{i+1}(x)/c_{i+1}, \text{ for } i = 0, 1, 2, \cdots$$

It is obvious that the division of polynomials will be done for less than n-1 times. If it is computed for n-1 times, then $\deg(q_i(x)) = 1$, and the leading coefficient of $r_i(x), q_i(x)$ is 1. So the division of polynomials involves $2\sum_{i=1}^{n-1} i = n^2 + \mathcal{O}(n)$ flops. If it has been computed for less than n-1 times, although $\deg(q_i(x)) > 1$, the times of polynomial division need to be reduced. It can be deduced that this part involves less than n^2 flops.

Second we discuss the computational complexity on the multiplication of polynomials.

$$u_{i+1}(x) \leftarrow [u_{i-1}(x) - q_i(x)u_i(x)]/c_{i+1}, \quad i = 0, 1, 2, \cdots$$

From Algorithm 4.1, we know that the multiplication of polynomials will be done for less than n-1 times. If it has been done for n-1 times, since $\deg(q_i(x)) = 1$, and the leading coefficient

of $r_i(x), q_i(x)$ is 1, the multiplication of polynomials requires $2\sum_{i=1}^{n-1} i = n^2 + \mathcal{O}(n)$ flops. If it has been computed for less than n-1 times, the multiple of division involves less than n^2 flops. So, all the amount of work is $2n^2$ flops. Thus, when $\deg(f(x)) = m$, the algorithm involves 3nm flops. So when $m \ll n$, we only require $\mathcal{O}(n)$ flops, which indicates that the algorithm reduces the computational complexity greatly.

With the same reason, when the scaled factor circulant matrix is singular, the method requires $4n^2 + \mathcal{O}(n)$ flops to compute the group inverse or the Moore-Penrose inverse.

6. Numerical Examples

Example 6.1. Let

$$A = \left(\begin{array}{rrrrr} 1 & 3 & 2 & 8 \\ 16 & 1 & 6 & 8 \\ 8 & 8 & 1 & 12 \\ 6 & 2 & 4 & 1 \end{array}\right).$$

Is the matrix A singular or nonsingular? If A is nonsingular, find the inverse of A.

Since $A = \text{scacirc}_{\mathcal{R}}(1,3,2,8) = I + 3\mathcal{R} + \mathcal{R}^2 + \mathcal{R}^3 = f(\mathcal{R})$, where $f(x) = 1 + 3x + x^2 + x^3$ is the representor of A, and

$$\mathcal{R} = \left(\begin{array}{rrrr} 0 & 1 & 0 & 0 \\ 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & 4 \\ 2 & 0 & 0 & 0 \end{array}\right),$$

it is known that A is a scaled factor circulant matrix. Let

$$r_{-1}(x) = g(x) = x^{4} - 16,$$

$$r_{0}(x) = f(x) = 1 + 3x + x^{2} + x^{3},$$

$$u_{-1}(x) = 0, \quad u_{0}(x) = 1.$$

Using Algorithm 4.1, we have

$$q_0(x) = x - 1, \quad r_1(x) = x^2 - x + \frac{15}{2}, \quad c_1 = -2, \quad u_1(x) = 0.5x - 0.5;$$

$$q_1(x) = x + 2, \quad r_2(x) = x + \frac{28}{5}, \quad c_2 = -2.5, \quad u_2(x) = 0.2x^2 + 0.2x - 0.8;$$

$$q_2(x) = x - \frac{33}{5}, \quad r_3(x) = 1, \quad c_3 = \frac{2223}{50},$$

$$u_3(x) = -\frac{289}{2223} + \frac{131}{2223}x + \frac{56}{2223}x^2 - \frac{10}{2223}x^3.$$

Since $r_3(x) = 1$, A is a nonsingular matrix and $u_3(x)$ is the representor of A^{-1} . Then

$$A^{-1} = \operatorname{scacirc}_{\mathcal{R}}\left(-\frac{289}{2223}, \frac{131}{2223}, \frac{112}{2223}, -\frac{80}{2223}\right)$$
$$= \frac{1}{2223} \begin{pmatrix} -289 & 131 & 112 & -80\\ -160 & -289 & 262 & 448\\ 448 & -80 & -289 & 524\\ 262 & 112 & -40 & -289 \end{pmatrix}$$

A New Algorithm for Computing the Inverse of the Scaled Factor Circulant Matrix

Example 6.2. Let

$$A = \begin{pmatrix} -4 & -3 & 2\\ 64 & -4 & -6\\ -96 & 32 & -4 \end{pmatrix}.$$

Is the matrix A singular or nonsingular? If A is singular, find the group inverse of A.

Since $A = \text{scacirc}_{\mathcal{R}}(-4, -3, 2) = -4I - 3\mathcal{R} + \mathcal{R}^2 = f(\mathcal{R})$, where $f(x) = -4 - 3x + x^2$ is the representor of A, and

$$\mathcal{R} = \left(\begin{array}{ccc} 0 & 1 & 0 \\ 0 & 0 & 2 \\ 32 & 0 & 0 \end{array} \right),$$

it is known that A is a scaled factor circulant matrix. Let $r'_{-1}(x) = g(x) = x^3 - 64$, $r'_0(x) = f(x) = -4 - 3x + x^2$. Using Algorithm 4.1, we have

$$q'_0(x) = x + 3$$
, $r'_1(x) = x - 4$, $c'_1 = 13$, $q'_1(x) = x + 1$, $r'_2(x) = 0$.

Then the largest common factor of f(x) and g(x) is $r'_1(x) = x - 4$. Consequently, A is a singular matrix. Let

$$r(x) = x - 4, \quad r_{-1}(x) = \frac{g(x)}{x - 4} = x^2 + 4x + 16,$$

$$r_0(x) = f(x)(x - 4) \mod(x^2 + 4x + 16) = 36x + 192, \quad u_{-1}(x) = 0, \quad u_0(x) = 1.$$

Using Algorithm 4.1, we have

$$q_0(x) = \frac{1}{36}x - \frac{1}{27}, \quad r_1(x) = 1, \quad c_1 = \frac{208}{9}, \quad u_1(x) = -\frac{1}{832}x + \frac{1}{624}.$$

Consequently,

$$u_1(x)r(x) = -\frac{1}{156} + \frac{1}{156}x - \frac{1}{832}x^2$$

is the representor of $A^{\#}$. Then

$$A^{\#} = \operatorname{scacirc}_{\mathcal{R}}\left(-\frac{1}{156}, \frac{1}{156}, -\frac{1}{416}\right) = \left(\begin{array}{ccc} -\frac{1}{156} & \frac{1}{156} & -\frac{1}{416} \\ -\frac{1}{13} & -\frac{1}{156} & \frac{2}{156} \\ \frac{32}{156} & -\frac{1}{26} & -\frac{1}{156} \end{array}\right).$$

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