# A COLUMN RECURRENCE ALGORITHM FOR SOLVING LINEAR LEAST SQUARES PROBLEM ${ }^{* 1)}$ 

J.X. Zhao<br>(Department of Mathematics, Nanjing University, Nanjing, China)


#### Abstract

A new column recurrence algorithm based on the classical Greville method and modified Huang update is proposed for computing generalized inverse matrix and least squares solution. The numerical results have shown the high efficiency and stability of the algorithm.


## 1. Introduction

Numerical method of a generalized inverse matrix and corresponding with the linear least squares is a standard tool for solving such problems as control, state evaluation and identification. Let A be an $m \times n$ real matrix. A real $n \times m$ matrix G is called the M-P generalized inverse matrix of $A$ if $G$ satisfies the following conditions:

$$
\begin{align*}
\text { (I) } A G A & =A, \quad(\mathrm{II}) \quad G A G=G \\
\text { (III) }(A G)^{T} & =A G, \quad(\mathrm{IV}) \quad(G A)^{T}=G A . \tag{1}
\end{align*}
$$

Usually, we write G

$$
A^{+}=G .
$$

The linear least square problem is defined as the minimization of the norm of the residual vector

$$
\begin{equation*}
\min _{x}\|A x-b\|_{2}^{2} \tag{2}
\end{equation*}
$$

where $b$ is an m-vector and $x$ is an n-vector. Thus, the least square solution of the minimum norm of problem (2) is

$$
x=A^{+} b
$$

One of the major stability indices for computing generalized inverse matrix or linear least squares problem is

$$
\kappa(A)=\|A\|\left\|A^{+}\right\| .
$$

[^0]An ill-conditioned matrix A, i.e. matrix with large $\kappa$, is quite common in control and identification problem ${ }^{[5]}$. It is thus important to have computational procedures suitable for solving ill-conditioned problem. An excellent survey on linear least square has been given in $\mathrm{Björck}^{[1]}$.

Since the Greville scheme ${ }^{[6]}$ is relative simple and is called G-method, it is adopted frequently for computing generalized inverse matrix in some cases. Computational practice and theoretical analysis show, however, when $A$ is an ill-conditioned matrix, that the solution computed by G-method may bear no resemblance to the true solution. On the other hand, modified Huang method, one of the ABS class, may be more stable than that of some classical matrix factorization method ${ }^{[2]}$. But this method is only fit for solving the problem where $m \leq n$.

Our aim of this paper is to describe a new modification of the classical Greville method, which retains the main advantages of the classical scheme but in many cases is more stable.

Throughout this paper, let $\|\cdot\|$ stand for the 2-norm of a matrix or a vector.

## 2. Greville Method and Its Modification

Let A be a matrix of order $m \times n$ and will be denoted by

$$
A=\left[a_{1}, a_{2}, \cdots, a_{n}\right] \in R^{m \times n}
$$

where $a_{i} \in R^{n}$ and $m \geq n$. By convention, we assume $\operatorname{rank}(A)=n$.
Denoted by

$$
\begin{equation*}
A_{1}=\left[a_{1}\right], \quad A_{k}=\left[A_{k-1}, a_{k}\right] \in R^{m \times k}, \quad k \leq n \tag{3}
\end{equation*}
$$

We have known that G-method is an electable method for computing generalized inverse matrix if A is not too ill-conditioned. The G-method proceeds as follows ${ }^{[6]}$.

## G-method:

Set

$$
A_{1}^{+}=a_{1}^{T} /\left(a_{1}^{T} a_{1}\right)
$$

For $\mathrm{k}=2: \mathrm{n}$ Compute

$$
d_{k}=A_{k-1}^{+} a_{k}, \quad c_{k}=a_{k}-A_{k-1} d_{k}
$$

Take

$$
y_{k}^{T}= \begin{cases}c_{k}^{+} & c_{k} \neq 0 \\ \left(1+d_{k}^{T} d_{k}\right) d_{k}^{T} A_{k-1}^{+} & c_{k}=0\end{cases}
$$

Compute

$$
A_{k}^{+}=\left[\begin{array}{c}
A_{k-1}^{+}-d_{k} y_{k}^{T} \\
y_{k}^{T}
\end{array}\right]
$$

Set

$$
A^{+}=A_{n}^{+}
$$

This algorithmic scheme is known as the Greville recurrence procedure. It is easy to see that the scheme is very compact and the numerical result is acceptable to well-behavior matrix.

Denote by $R(A)$ and $R(A)^{\perp}$ the range of A and corresponding orthogonal complement, respectively. Since

$$
\begin{align*}
c_{k} & =a_{k}-A_{k-1} d_{k} \\
& =\left(I_{m}-A_{k-1} A_{k-1}^{+}\right) a_{k} \tag{4}
\end{align*}
$$

$c_{k}$ is an orthogonal projection vector of $a_{k}$ onto $R\left(A_{k-1}\right)^{\perp}$. The following lemma thus can be obtained directly[6].

Lemma. The orthogonal projection $c_{k}=0$ in $G$-method iff $a_{k}$ is a linear combination of the vector system $a_{1}, a_{2}, \cdots, a_{k-1}$, i.e.

$$
a_{k} \in R\left(A_{k-1}\right) .
$$

In [4] the relation between the G-method and classical Gram-schmidt process is analyzed and the G-method for computing M-P inverse matrix, when $\operatorname{rank}(A)=n \leq m$, is entirely equivalent to Gram-schmidt process from a purely theoretical viewpoint. In other words, we can prove the following relations to G-method[4].

Theorem. If $c_{1}, c_{2}, \cdots, c_{n}$ are obtained by $G$-method, then

1. $c_{k}^{T} a_{j}=c_{k}^{T} c_{j}=0, \quad k=1,2, \cdots, n ; \quad j=1,2, \cdots, k-1 ;$
2. $\quad c_{i}^{T} a_{i}=c_{i}^{T} c_{i}, \quad i=1,2, \cdots, n$;
3. $\left(I_{m}-\Sigma_{l=1}^{k-1} c_{l} c_{l}^{+}\right) a_{j}=\left(I_{m}-\sum_{l=1}^{k-1} c_{l} c_{l}^{+}\right) c_{j}=0, \quad j=1,2, \cdots, k$.

Proof. By induction.
From the above discussion, we have known that G-method is equivalent to classical Gram-schmidt process. Therefore, its numerical stability is relatively weak for illconditioned problem. In general, the modified Gram-schmidt process is more stable than that of classical scheme. Based on the analogous idea, we can obtain the following modified Greville method.

Take

$$
A_{1}^{+}=a_{1}^{+},
$$

Compute

$$
d_{2}^{(j)}=A_{1}^{+} a_{j}, \quad j=2,3, \cdots, n .
$$

Let

$$
d_{2}=d_{2}^{(2)} .
$$

Then using the relations

$$
d_{3}^{(j)}=A_{2}^{+} a_{j}=\left[\begin{array}{c}
A_{1}^{+}-d_{2} y_{2}^{T} \\
y_{2}^{T}
\end{array}\right] a_{j}
$$

$$
=\left[\begin{array}{c}
d_{2}^{(j)}-\left(y_{2}^{T} a_{j}\right) d_{2} \\
y_{2}^{T} a_{j}
\end{array}\right],
$$

we can compute $d_{3}^{(j)}, j=3,4, \cdots, n$.
Set

$$
d_{3}=d_{3}^{(3)}=A_{2}^{+} a_{3} .
$$

At the beginning of k -th $(k \leq n)$ step, assume that $d_{k-1}^{(j)}, j=k, k+1, \cdots, n$, have already been computed. Then compute

$$
\begin{align*}
d_{k}^{(j)} & =A_{k-1}^{+} a_{j}=\left[\begin{array}{c}
d_{k-1}^{(j)}-\left(y_{k-1}^{T} a_{j}\right) d_{k-1} \\
y_{k-1}^{T} a_{j}
\end{array}\right],  \tag{5}\\
j & =k, k+1, \cdots, n .
\end{align*}
$$

Set

$$
d_{k}=d_{k}^{(k)}=A_{k-1}^{+} a_{k} .
$$

We can see that the vector system $c_{1}, c_{2}, \cdots, c_{k}, k \leq n$, obtained by G-method must be in $R\left(A_{k}\right)$. Therefore,

$$
\begin{equation*}
A_{k} A_{k}^{+} c_{i}=c_{i} \quad i=1,2, \cdots, k \tag{6}
\end{equation*}
$$

This means that the eigenvectors of $A_{k} A_{k}^{+}$associated with eigenvalue 1 are $c_{i}, i=$ $1,2, \cdots, k$. It follows from the spectral decomposition that

$$
A_{k} A_{k}^{+}=c_{1} c_{1}^{+}+c_{2} c_{2}^{+}+\cdots+c_{k} c_{k}^{+}
$$

Let

$$
\begin{align*}
c_{k}^{(j)} & =\left(I_{m}-A_{k} A_{k-1}^{+}\right) a_{j} \\
& =\left(I_{m}-\Sigma_{l=1}^{k-1} c_{l} c_{l}^{+}\right) a_{j} \\
& =c_{k-1}^{(j)}-c_{k-1} c_{k-1}^{+} a_{j} \\
& =c_{k-1}^{(j)}-\frac{\left(c_{k-1}^{T} a_{j}\right) c_{k-1}}{c_{k-1}^{T} c_{k-1}} \quad j=k, k+1, \cdots, n . \tag{7}
\end{align*}
$$

Take

$$
c_{k}=c_{k}^{(k)}
$$

Based on the above discussion, we can lead the modified Greville method, called the MG-method, which is equivalent to modified Gram-schmidt process. To ensure the numerical stability of G-method, a pivoting strategy is often performed in the size of $\left\|c_{k}\right\|$.

## MG-method:

Compute

For $k=2: n$,

1. Compute

$$
d_{k}^{(j)}=\left[\begin{array}{c}
d_{k-1}^{(j)}-\left(y_{k-1}^{T} a_{j}\right) d_{k-1} \\
y_{k-1}^{T} a_{j}
\end{array}\right], j=k, k+1, \cdots, n
$$

2. Compute

$$
c_{k}^{(j)}=c_{k-1}^{(j)}-\frac{\left(c_{k-1}^{T} a_{j}\right) c_{k-1}}{c_{k-1}^{T} c_{k-1}}, j=k, k+1, \cdots, n .
$$

Pivoting strategies:

$$
c_{\max }=\left\|c_{k}^{\left(j_{0}\right)}\right\|=\max _{k \leq j \leq n}\left(\left\|c_{k}^{(j)}\right\|\right)
$$

Interchanging:

$$
a_{j_{0}} \leftrightarrow a_{k}, \quad c_{k}^{\left(j_{0}\right)} \leftrightarrow c_{k}^{(k)}, \quad d_{k}^{\left(j_{0}\right)} \leftrightarrow d_{k}^{(k)}
$$

Let

$$
c_{k}=c_{k}^{(k)}, \quad d_{k}=d_{k}^{(k)}
$$

3. Take

$$
y_{k}^{T}=c_{k}^{+} .
$$

4. Compute

$$
A_{k}^{+}=\left[\begin{array}{c}
A_{k-1}^{+}-d_{k} y_{k}^{T} \\
y_{k}^{T}
\end{array}\right] .
$$

Set

$$
A^{+}=A_{n}^{+}
$$

Furthermore, the numerical experiments show that the updating precision of $c_{k}, k=$ $1,2, \cdots, n$, makes great influence on the numerical stability in the G-method. On the other hand, the ABS methods for proceeding successively to orthogonal bases of $R(A)$ have been obtained for the case where $m \leq n[2]$.

## ABS-updating:

Let $H_{1}$ be an arbitrary $n \times n$ nonsingular matrix. For $\mathrm{i}=1,2, \cdots, \mathrm{n}$, compute

$$
\begin{equation*}
p_{i}=H_{i} z_{i}, \tag{i}
\end{equation*}
$$

where $z_{i} \in R^{n}$ is arbitrary, subject to $z_{i}^{T} p_{i} \neq 0$;
(ii) $\quad H_{i+1}=H_{i}-H_{i} w_{i} a_{i}^{T} H_{i} /\left(a_{i}^{T} H_{i} w_{i}\right)$,
where $w_{i} \in R^{n}$ is arbitrary, subject to $a_{i}^{T} H_{i} w_{i} \neq 0$.
The vector system $p_{1}, p_{2}, \cdots, p_{k}$ obtained by this algorithm satisfy

$$
p_{k}^{T} a_{j}=0, j=1,2, \cdots, k-1 .
$$

Among the particular algorithms in the ABS class obtained by making specific choices of the parameters, the modified Huang method is of special interest:

$$
H_{1}=I_{n} .
$$

For $i=1,2, \cdots, n$. do

$$
\begin{aligned}
z_{i} & =H_{i} a_{i}, w_{i}=z_{i} / z_{i}^{T} z_{i} \\
p_{i} & =H_{i}^{T} z_{i}, \\
H_{i+1} & =H_{i}-H_{i} a_{i} w_{i}^{T} H_{i} .
\end{aligned}
$$

This algorithm is numerically more stable than that of the Huang method or its analogues for the determined or underdetermined problem[2].

Here we wish to extend the modified Huang updating to overdetermined problem. It is to say that the orthogonal projection vector system $c_{1}, c_{2}, \cdots, c_{n}$, can be obtained by the modified Huang updating in the MG-method. Based on this idea, we can comprise a numerical method, called MHG-method, for computing the M-P generalized inverse.

MHG-method:

1. Set $H_{1}=I_{m}$. and compute $A_{1}^{+}$,
2. For $k=2: n$ compute
2.1

$$
\begin{aligned}
d_{k}^{(j)} & =\left[\begin{array}{c}
d_{k-1}^{(j)}-\left(y_{k}^{T} a_{j}\right) d_{k-1} \\
y_{k}^{T} a_{j}
\end{array}\right], j=k, k+1, \cdots, n \\
d_{k} & =d_{k}^{(k)}
\end{aligned}
$$

2.2

$$
\begin{aligned}
z_{k} & =H_{k} a_{k}, \quad w_{k}=z_{k}^{T} z_{k}, \\
c_{k} & =H_{k}^{T} z_{k}, \\
H_{k+1} & =H_{k}-H_{k} a_{k} w_{k}^{T} H_{k} .
\end{aligned}
$$

2.3

$$
A_{k}^{+}=\left[\begin{array}{c}
A_{k-1}^{+}-d_{k} c_{k}^{+} \\
c_{k}^{+}
\end{array}\right]
$$

3. Set

$$
A^{+}=A_{n}^{+}
$$

In many cases, above the algorithm is much stabler than that of the G-method or MG-method for ill-conditioned problems.

## 3. Solution of Linear Least Squares Problem

Applying the above given MHG-method for computing the generalized inverse, we here can directly extend to solving overdetermined linear least squares problem

$$
\begin{equation*}
\min _{x}\|A x-b\|^{2} \tag{8}
\end{equation*}
$$

where A is an $m \times n$ matrix, $m \geq n$, and b is an $m$-vector.
Assume that $A_{k-1}^{+}$has been computed at the beginning of k-th step. Then we can form

$$
x_{k-1}=A_{k-1}^{+} b
$$

and

$$
A_{k}=\left[A_{k-1}, a_{k}\right] .
$$

Therefore, we have

$$
\begin{aligned}
x_{k} & =A_{k}^{+} b=\left[\begin{array}{c}
A_{k-1}^{+} b-\left(y_{k}^{T} b\right) d_{k} \\
y_{k}^{T} b
\end{array}\right] \\
& =\left[\begin{array}{c}
x_{k-1}-x_{k}^{\prime} d_{k} \\
x_{k}^{\prime}
\end{array}\right]
\end{aligned}
$$

where

$$
x_{k}^{\prime}=y_{k}^{T} b
$$

and $y_{k}, d_{k}$ are computed by the MHG-method.
To be more compact, the new algorithm takes the vector b in the right hand side of the above equality as $(\mathrm{n}+1)$ th column of A . Based on the above idea, the following algorithm has been given, called the MHGS-method, for solving the linear least squares problem.

## MHGS-method:

1. Set $[A: b]=\left[a_{1}, a_{2}, \cdots, a_{n+1}\right] \in R^{m \times(n+1)}, H_{1}=I_{m}, R=\left[r_{i j}\right]=0 \in R^{n \times(n+1)}$, a working array of the upper triangular matrix;
2. For $i=1,2, \cdots, n$ do

$$
\begin{aligned}
& z_{i}=H_{i} a_{i} \\
& w_{i}=z_{i} / z_{i}^{T} z_{i}, \\
& c_{i}=H_{i}^{T} z_{i}
\end{aligned}
$$

If $\mathrm{i}=1$ then $r_{1 j}=c_{1}^{T} a_{j} / c_{1}^{T} c_{1}, \quad j=2, \cdots, n+1$;
else for $j=i+1, \cdots, n+1$ do
for $k=1,2, \cdots, i-1$ do
$r_{k j}=r_{k j}-\frac{c_{i}^{T} a_{j}}{c_{i}^{T} c_{i}} r_{k i}$,
$r_{i j}=\frac{c_{i}^{T} a_{i}}{c_{i}^{T} c_{i}}$.
$H_{i+1}=H_{i}-H_{i} a_{i} w_{i}^{T} H_{i}$.
3. $x_{i}=r_{i . n+1}, i=1,2, \cdots, n$.
$x=\left(x_{1}, x_{2}, \cdots, x_{n}\right)^{T}$ is the least squares solution of the minimum norm of (8).

## 4. Numerical Experiments

The algorithm described in previous sections has been implemented and has been run in double precision on VAX computer of the University of Trento. The considered linear least squares problems,

$$
\min _{x}\|A x-b\|^{2}
$$

have the following data.
Problem 1. (ill-conditioned problem)

$$
\begin{aligned}
A & =\left[a_{i j}\right]=[1.0 /(i+j-1.0)] \in R^{m \times n} \\
b & =\left(b_{1}, b_{2}, \cdots, b_{m}\right)^{T} \in R^{m} \\
b_{i} & =\Sigma_{j=1}^{n} a_{i j}, \quad i=1,2, \cdots, m
\end{aligned}
$$

Problem 2. (mid-conditioned problem)

$$
\begin{aligned}
A & =\left[a_{i j}\right]=[\max (i, j)] \in R^{m \times n} \\
b & =\left(b_{1}, b_{2}, \cdots, b_{m}\right)^{T} \in R^{m}, \\
b_{i} & =\Sigma_{j=1}^{n} a_{i j}, \quad i=1,2, \cdots, m .
\end{aligned}
$$

Problem 3. (mid-conditioned problem)

$$
\begin{aligned}
A & =\left[\begin{array}{cccccc}
n & n-1 & n-2 & \ldots & 2 & 1 \\
n-1 & n-1 & n-2 & \ldots & 2 & 1 \\
n-2 & n-2 & n-2 & \ldots & 2 & 1 \\
\ldots & \ldots & \ldots & \ldots & \ldots & \ldots \\
2 & 2 & 2 & \ldots & 2 & 1 \\
1 & 1 & 1 & \ldots & 1 & 1
\end{array}\right]=\left[a_{i j}\right] \in R^{n \times n}, \\
b & =\left(b_{1}, b_{2}, \cdots, b_{n}\right)^{T} \in R^{n}, \\
b_{i} & =\sum_{j=1}^{n} a_{i j}, \quad i=1,2, \cdots, n .
\end{aligned}
$$

## Algorithm:

G-method: Column Greville method, CGLSP ${ }^{2)}$-method : Algorithm given in [5],

[^1]MHGS-method: Algorithm obtained in this paper.

The following tables present the computed precision with three different algorithms for the given problems. These results indicate that the new method (MHGS-method) obtained in this paper is fairly efficient, especially as an ill-conditioned least square as problem 1. The index of the relative accuracy is defined by

$$
P=\frac{\left\|x-x^{*}\right\|_{2}}{\left\|x^{*}\right\|_{2}}
$$

where $x^{*}$ is the true solution and x is a computed solution.
The condition numbers of the given matrices, cond $(\mathrm{A})$, are computed by MATLAB.

Table 1. The relative accuracy P of Problem $1(m=n)$.

| m | Cond(A) | G-method | CGLSP-method | MHGS-method |
| :--- | :---: | :---: | :---: | :---: |
| 5 | $4.7661 \mathrm{E}+05$ | $1.1666822 \mathrm{E}-07$ | $2.9090133 \mathrm{E}-05$ | $2.1568097 \mathrm{E}-12$ |
| 10 | $1.6025 \mathrm{E}+13$ | $>1$ | $3.7835974 \mathrm{E}-04$ | $6.1374327 \mathrm{E}-09$ |
| 15 | $3.3634 \mathrm{E}+17$ | $>1$ | $1.1718496 \mathrm{E}-03$ | $7.3047523 \mathrm{E}-09$ |
| 20 | $4.0020 \mathrm{E}+19$ | $\gg 1$ | $5.2199944 \mathrm{E}-04$ | $2.4599253 \mathrm{E}-08$ |
| 25 | $1.7138 \mathrm{E}+18$ | $\gg 1$ | $8.5164341 \mathrm{E}-04$ | $1.0516242 \mathrm{E}-08$ |
| 30 | $4.5397 \mathrm{E}+18$ | $\gg 1$ | $1.2079764 \mathrm{E}-03$ | $2.2723464 \mathrm{E}-08$ |
| 35 | $5.8601 \mathrm{E}+18$ | $\gg 1$ | $1.5771304 \mathrm{E}-03$ | $2.0508478 \mathrm{E}-08$ |
| 40 | $7.2654 \mathrm{E}+18$ | $\gg 1$ | $6.1066032 \mathrm{E}-04$ | $5.0091549 \mathrm{E}-08$ |

Table 2. The relative accuracy P of Problem $2(m=n)$.

| m | Cond(A) | G-method | CGLSP-method | MHGS-method |
| :---: | :--- | :--- | :--- | :--- |
| 5 | $7.0766 \mathrm{E}+01$ | $7.3474726 \mathrm{E}-14$ | $4.7808928 \mathrm{E}-16$ | $2.5225527 \mathrm{E}-16$ |
| 10 | $2.8919 \mathrm{E}+02$ | $1.7468340 \mathrm{E}-11$ | $9.9344994 \mathrm{E}-16$ | $3.2823535 \mathrm{E}-15$ |
| 15 | $6.4639 \mathrm{E}+02$ | $1.8298022 \mathrm{E}-07$ | $1.4754814 \mathrm{E}-15$ | $6.2574871 \mathrm{E}-15$ |
| 20 | $1.1425 \mathrm{E}+03$ | $3.3819112 \mathrm{E}-04$ | $4.3725890 \mathrm{E}-15$ | $1.5046502 \mathrm{E}-14$ |
| 25 | $1.7775 \mathrm{E}+03$ | $0.2837466 \mathrm{E}+00$ | $5.6821201 \mathrm{E}-15$ | $1.9495403 \mathrm{E}-14$ |
| 30 | $2.5515 \mathrm{E}+03$ | $>1$ | $9.2010109 \mathrm{E}-15$ | $2.2474395 \mathrm{E}-14$ |
| 35 | $3.4644 \mathrm{E}+03$ | $>1$ | $1.1894571 \mathrm{E}-14$ | $4.6867962 \mathrm{E}-14$ |
| 40 | $4.5163 \mathrm{E}+03$ | $>1$ | $1.6454910 \mathrm{E}-14$ | $5.3042908 \mathrm{E}-14$ |

Table 3. The relative accuracy P of Problem $3(m=n)$.

| m | Cond(A) | G-method | CGLSP-method | MHGS-method |
| :---: | :--- | :--- | :--- | :--- |
| 5 | $4.5455 \mathrm{E}+01$ | $0.0000000 \mathrm{E}+00$ | $1.4057041 \mathrm{E}-16$ | $0.0000000 \mathrm{E}+00$ |
| 10 | $1.7508 \mathrm{E}+02$ | $3.6570011 \mathrm{E}-16$ | $3.3419257 \mathrm{E}-16$ | $0.0000000 \mathrm{E}+00$ |
| 15 | $3.8582 \mathrm{E}+02$ | $4.1579343 \mathrm{E}-15$ | $6.8814798 \mathrm{E}-16$ | $0.0000000 \mathrm{E}+00$ |
| 20 | $6.7762 \mathrm{E}+02$ | $1.4859172 \mathrm{E}-14$ | $1.6301262 \mathrm{E}-15$ | $0.0000000 \mathrm{E}+00$ |
| 25 | $1.0505 \mathrm{E}+03$ | $1.6141173 \mathrm{E}-14$ | $3.4069208 \mathrm{E}-15$ | $0.0000000 \mathrm{E}+00$ |
| 30 | $1.5044 \mathrm{E}+03$ | $3.4304422 \mathrm{E}-14$ | $3.9338318 \mathrm{E}-15$ | $0.0000000 \mathrm{E}+00$ |
| 35 | $2.0394 \mathrm{E}+03$ | $2.7719791 \mathrm{E}-14$ | $8.5910155 \mathrm{E}-15$ | $0.0000000 \mathrm{E}+00$ |
| 40 | $2.6554 \mathrm{E}+03$ | $4.8073519 \mathrm{E}-14$ | $9.9144612 \mathrm{E}-15$ | $0.0000000 \mathrm{E}+00$ |

Table 4. The relative accuracy of the MHGS-method
for the ill-conditioned problem $1(m>n)$.

| $m$ | $n$ | Cond(A) | relative accuracy P |
| :---: | :---: | :--- | :--- |
| 150 | 100 | $1.5173 \mathrm{E}+18$ | $3.3504126 \mathrm{E}-08$ |
| 150 | 110 | $2.1010 \mathrm{E}+18$ | $4.0557843 \mathrm{E}-08$ |
| 150 | 120 | $2.5797 \mathrm{E}+18$ | $4.6187279 \mathrm{E}-08$ |
| 150 | 130 | $4.0761 \mathrm{E}+18$ | $5.2436966 \mathrm{E}-08$ |
| 150 | 140 | $1.0964 \mathrm{E}+19$ | $9.6172765 \mathrm{E}-08$ |
| 150 | 150 | $1.9002 \mathrm{E}+20$ | $2.0729776 \mathrm{E}-07$ |
| 200 | 150 | $2.7419 \mathrm{E}+18$ | $4.8961957 \mathrm{E}-08$ |
| 500 | 10 | $9.9475 \mathrm{E}+09$ | $1.6412854 \mathrm{E}-09$ |
| 500 | 100 | $6.1819 \mathrm{E}+17$ | $3.7023077 \mathrm{E}-08$ |

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[^0]:    * Received May 23, 1994.
    ${ }^{1)}$ On leave at the Department of Mathematics, University of Trento, 38050 POVO (Trento), Italy.

[^1]:    ${ }^{2)}$ Modified conjugate gradient method for least squares problem

